MOV DWORD PTR DS: LECX+60 MOV ECX:DWORD PTR SS:LEC MOV ECX:DWORD PTR SS:LEC MOV ECX:DWORD PTR SS:LEC MOV ECX:DWORD PTR SS:LECX+60 MOV ECX

> MOV DWORD PTR DS: CEDX MOV ECX, DWORD PTR SS

Binary Obfuscation from the Top Down

How to make your compiler do your dirty work.

MOV CL. BY

JNZ SHO

Why Top Down?

Assembly, while "simple," is tedious.
It's easier for us to write higher-level code.

- Some of us.
- Why do it by hand when you can be lazy?

LEA EDX, DWORD MOV CL, BYT

JNZ

SH

What's the purpose of obfuscation?

MOV EST DWORD

CL.BY

JNZ

SH

- To waste time.
- To intimidate.
- To be a total jerk.

What tools will be used?

- C and C++
- MSVC++ for compilation (sorry)

MOV EST DWO

MOV CL. BY

JNZ

SH

What will not be covered?

- Anti-debug
- Source obfuscation where it does not relate to binary transformations
- Obfuscation effectiveness
- Post-compilation obfuscation

Important Basics

Hopefully we can get through this really quickly.

JNZ

SHO

CALL 00401028 POP ECX POP ECX PUSH 195 PUSH 00417FF8 MOV ESI,00417E60 Fun With Pointers

car cdr cdar cdadr cdddr caar caaaar caaaaar caaaaaaar

Function Pointers

- Like string-format vulnerabilities, function pointers are ancient Voodoo.
- I honestly don't know who thought these were a good idea, but I freakin' love 'em.
- See src/funcptr.c

Function Pointers

int foo (void) {
 return 949;

int bar (void) {
 int (*fooPtr)(void);
 fooPtr = foo;
 return fooPtr();

PUSH 00417FF8 MOV ESI,00417E60

Method Pointers

- Abuse of method pointers would probably make Bjarne Stroustrup really angry.
- There is also one thing uglier than function pointers. That's method pointers.
- See src/methodptr.cpp

Method Pointers

int MyClass::foo(void) {
 return 310;

int bar (void) { MyClass baz; int (MyClass:*fooPtr)(void); fooPtr = & MyClass::foo; return (MyClass.*baz)fooPtr();

PUSH 00417FF8 MOV ESI,00417E60 **Calling Conventions**

I really want to write a clever pun about payphones and DEFCON, but I just can't.

JNZ

SH

Calling Conventions

 When making a function call, there are a few ways to do it:

ESTIDWO

CALL

LEA EDX, DWORD MOV CL, BYT

stdcall
cdecl
fastcall
thiscall

Calling Conventions

- stdcall
 - Push arguments onto stack
 - Called function pops from stack

CAL

• Cleans up its own mess.

Calling Conventions

• cdecl

Push arguments onto stack
Called function pops from stack

• Called function cleans up the mess

LEA EDX, DWORD MOV CL, BY

Calling Conventions

- fastcall
 - First two arguments less than a DWORD moved into ecx and edx respectively
 - Rest are pushed onto the stack
 - Called function pops from the stack
 - Called function cleans up the mess

Calling Conventions

- thiscall
 - Used when a function within a class object is called
 - "this" pointer moved into ecx
 - Function arguments pushed onto stack
 - Called function pops from stack
 - Cleans up its own mess

Compiler Optimizations

The Dragon Book: Not Just for Furries Anymore

CL.BY

JNZ

SH

SUB EAX, EDX PUSH EAX PUSH ESI OO401028 POP ECX POP ECX PUSH 195 PUSH 00417FF8 MOV ESI,00417E60

Compiler Optimizations

- Control-flow analysis
- Variable analysis
- Reach-of-use
- The volatile keyword

LEA EDX. DWOR

CAL

Compiler Optimizations

- At compile time, your code is separated into multiple blocks.
- A "block" consists of code separated by conditional (e.g. JLE, JNE, etc.) and unconditional jumps (e.g. CALL and JMP).
- How this code is organized and how the jumps occur affects the optimization of the program.

Compiler Optimizations MOV EAX,949 XOR EAX,310 CMP EAX,0 JNE z0r

z0r: XOR EAX,310 PUSH EAX XOR EAX,949 LEAVE RETN

PUSH 00417FF8 MOV ESI,00417E60

Compiler Optimizations

MOV EAX,949 XOR EAX,310 CMP EAX,0 JNE z0r

lol lemme fix this

z0r: XOR EAX,310 PUSH EAX XOR EAX,949 LEAVE RETN

PUSH 00417FF8 MOV ESI,00417E60

Compiler Optimizations

MOV EAX,949 XOR EAX,310 XOR EAX,310 PUSH EAX

, 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310

Compiler Optimizations

- The compiler also looks at your variables to make sure you're not doing anything repetitive or inconsequential.
- Algorithms like the directed acyclic graph (DAG) algorithm and static variable analysis make sure memory and math are fully optimized.

Compiler Optimizations

MOV EAX,949 XOR EAX,310 XOR EAX,310 PUSH EAX

, 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310 , 310

Compiler Optimizations

MOV EAX,949 XOR EAX,310 XOR EAX,310 PUSH EAX

lol seriously?

CAL

Compiler Optimizations

MOV EAX,949 PUSH EAX

JNZ

SH

Compiler Optimizations

MOV EAX,949 XOR EAX,310 CMP EAX,0 JNE z0r

z0r: XOR EAX,310 PUSH EAX

XOR EAX,949 LEAVE RETN

MOV EAX,949 PUSH EAX

PUSH EAX PUSH EAX PUSH EST POP ECX POP ECX PUSH 00417FF8 OV EST,00417FF8

Compiler Optimizations

- Your compiler is a neat-freak.
- If the compiler notices it doesn't need a variable anymore, it's just going to get rid of it, no matter what else you do to it.

CL.BY

Compiler Optimizations

MOV EAX,949 MOV EBX,310 MOV ECX,213 XOR EAX,EBX ADD EBX,EAX SUB EAX,EAX PUSH EBX PUSH EAX

Compiler Optimizations

MOV EAX,949 MOV EBX,310 MOV ECX,213 XOR EAX,EBX ADD EBX,EAX SUB EAX,EAX PUSH EBX PUSH EAX

Compiler Optimizations

MOV EAX,949 MOV EBX,310 XOR EAX,EBX ADD EBX,EAX SUB EAX,EAX PUSH EBX PUSH EAX

Compiler Optimizations

- There exist cases (mostly in hardware development) where you do NOT want your compiler to optimize your variable.
- This is where the volatile keyword comes in.
- Making your variable volatile tells the compiler not to do any optimizations to it.

Compiler Optimizations

volatile int foo; volatile char bar; volatile uint32 t baz;

CALL 00401028 PUSH 004012FA PUSH ESI PUSH 195 PUSH 00417FF8 MOV ESI,00417E60

Compiler Optimizations

int x; x = 7; x <<= 2; x *= 2; x *= 2; x -= 12; x -= 12; x += (x*x)<<2; printf("%d\n", x);

ALL 00401028 POP ECX POP ECX PUSH 195 PUSH 00417FF8 MOV ESI,00417E60

Compiler Optimizations

int x; x = 7; x <<= 2; x *= 2; x *= 2; x -= 12; x += (x*x)<<2; printf("%d\n", x);

PUSH 1E6C PUSH "%d\n" CALL \$PRINTF

PUSH EAX PUSH EST OO401028 POP ECX POP ECX PUSH 00417FF8 MOV ESI,00417E60

Compiler Optimizations

volatile int x; x = 7; x <<= 2; x *= 2; x *= 2; x -= 12; x += (x*x)<<2; printf("%d\n", x);

ALL 00401028 POP ECX POP ECX PUSH 00417FF8 MOV ESI,00417FF8

Compiler Optimizations

volatile int x; x = 7; x <<= 2; x *= 2; x *= 2; x -= 12; x += (x*x)<<2; printf("%d\n", x);

MOV [ESP],7 SHL [ESP],2 MOV EAX, [ESP] ADD EAX, EAX MOV [ESP], EAX ADD [ESP], -0CMOV ECX, [ESP] MOV EDX, [ESP] MOV EAX, [ESP] IMUL ECX, EDX

PUSH 00417FF8 MOV ESI,00417E60 **Binary Formats**

Everything is a file.

JNZ

SH

CALL

Binary Formats

- The most common formats you'll likely come across are the PE file format (Windows) and the ELF format (Linux).
- Both of these formats have a "table" they use for external library calls such as printf, execv, etc.
- For Windows it's called the IAT. For Linux it's the PLT.

Binary Formats

- If you obfuscate function pointers, they will likely not show up in those lists and therefore cause your library calls to fail.
- Circumventing this issue will be covered later.

Methods of Analysis

Know your opponent!

JNZ

SH

CALL

- Someone can easily figure out the gist of what your program is doing by analyzing any of the API calls you make.
- There exist a few programs out there that already do this for you: VirusTotal and ZeroWine.

Methods of Analysis

 VirusTotal (virustotal.com) is a website that allows you to upload suspected malware files and analyze them against over thirty different scanners.

 At the end of the analysis is a list of all recognized Windows API calls made by the program, as well as various data sections within.

Methods of Analysis

 ZeroWine (zerowine.sourceforge.net) is a malware analysis tool that executes a program in a controlled environment and collects data.

 This, too, collects and reports on API calls made by the program, as well as any possible servers it may have contacted or files it may have written.

- When analyzing a binary, there are two schools of analysis: live-code and dead-code.
- Dead-code is exactly how it sounds: you look at the binary, as-is, without executing.
- Live-code is the opposite: you run the program and watch what it does.

- VirusTotal employs dead-code analysis. It simply reads the binaries uploaded to it, scans it with various virus scanners and reports.
- ZeroWine, however, employs live-code analysis. It runs the suspected program in a controlled environment and watches what happens.

- Dead-code analysis can be frustrated through polymorphism.
- Live-code analysis can be frustrated through hiding, obfuscating and redirecting data and control-flow under the eyes of the reverser.

Obfuscation

We're almost at the fun part, I promise!

JNZ

SH

CALL

Obfuscation

• There are three separate classes of obfuscation.

SH

CALL

Layout
Control-flow
Data

Obfuscation

- Layout obfuscation essentially means scrambling the program around at the source-level.
- The International Obfuscated C Contest (ioccc.org) is a perfect example of this.

Obfuscation Anders Gavare, <u>http://www0.us.ioccc.org/2004/gavare.c</u>

X=1024; Y=768; A=3;

J=0;K=-10;L=-7;M=1296;N=36;O=255;P=9; =1<<15;E;S;C;D;F(b){E="1""111886:6:??AAF" "FHHMMOO55557799@@>>>BBBGGIIKK"[b]-64;C="C@=::C@@==@=:C@=:C@=:C0""31/513/5131/" "31/531/53"[b]-64;S=b<22?9:0;D=2;}I(x,Y,X){Y?(X^=Y,X*X>x?(X^=Y):0, I(x,Y/2,X) ,0);}p;q(c,x,y,z,k,l,m,a, $b) \{ F(c) \}$)):(E=X); $H(x) \{I(x,$;z-=C*M y*y/M+z);x-=E*M ;y-=S*M ;b=x* x/M+ p=((b=a*a/M-*z/M-D*D*M;a=-x *k/M -y*1/M-z*m/M; b)>=0?(I (b*M,____ a+(a>b -1.0); Z; W; O=E, ?-b:b)): ,0),b z,k,l, <44?(q(c,x ,y,z,k, (c,x,y, m,a){Z=! c? -1:Z;c ||Z<0) 1,m,0,0),(p> 0&&c!=)?(W= a&& (p<W p,Z=c):0,0(C+ 1, х,у,z, k,1, m,a)):0 ;}Q;T; b,V){0(0 h,i,j,d,a, ,e,f,q,h,i,j,a);d>0 U;u;v;w ;n(e,f,q, &&Z >= 0? (e+=h*W/M,f+=i*W/M,q+=j*W/M,F(Z),u=e-E*M,v=f-S*M,w=q-C*M,b=(-2*u-2*v+w) /3,H(u*u+v*v+w*w),b/=D,b*=b,b*=200,b/=(M*M),V=Z,E!=0?(u=-u*M/E,v=-v*M/E,w=-w*M/E) E):0,E=(h*u+i*v+j*w)/M,h-=u*E/(M/2),i-=v*E/(M/2),j-=w*E/(M/2),n(e,f,q,h,i,j,d-1) ,Z,O,O),Q/=2,T/=2, U/=2, V=V<22?7: (V<30?1:(V<38?2:(V<44?4:(V==44?6:3)))),Q+=V&1?b:0,T +=V&2?b :0,U+=V &4?b:0) :(d==P?(q+=2 250*U/M, T=255,j=q>0?q/8:q/ >0?(U= *j/M**,**Q =255-20):0,j j -150*U/M, U=255-100 *U/M):(U =i*i /M,U<M /5?(Q=255-210*U -=M/5, Q=213-110*U/M, U = 255/M,T=255-435*U -720* U/M):(U M, U = 111/M,T=168-113*U), d!=P?(Q/=2, T/=2)/ -85*U/M) ,U/=2):0);Q=Q< 0>0? O; T=T<0?0:T>O?O:T;U=U<0?0: 0?0: 0: U>O?O:U;}R;G;B *40*(A*x +a)/X/A-M*20, M*K, M;t(x,y b){n(M*J+M ,a, *L-M*30*(A*y+b)/Y/A+M*15,0,M,0,P, -1,0,0);R+=Q +=U;++a<A?t(x,y,a, ;G+=T;B b):(++b<A?t(x,y,0,b):0);}r(x,y){R=G=B=0;t(x,y,0,0);x<X?(printf("%c%c%c",R/A/A,G /A/A,B/A/A),r(x+1,y)):0;}s(y){r(0,--y?s(y),y:y);}main(){printf("P6\n%i %i\n255" "\n",X,Y);s(Y);}

UV ESI 00417E

Obfuscation

- Control-flow obfuscation involves twisting the typical downward-flow of a program to into spaghetti code.
- It has the added benefit of obfuscating source while simultaneously upsetting the normal flow a reverse-engineer is used to.

Obfuscation

- Data obfuscation involves masking whatever data you have in your program by any means.
- Strings, numbers, even functions within your program can be masked, obfuscated, interwoven or encrypted without handwriting any assembly.

Obfuscation Techniques

Now the fun begins.

JNZ

CALL

Obfuscation Techniques

- The goal is to obfuscate the binary without doing binary transformations.
- We know how the compiler optimizes, what it does to our data and how it stores some information important for programmatic logic.
- With this in mind, we can now leverage our code against the compiler.

POP ECX POP ECX PUSH 195 PUSH 00417FF8 MOV ESI,00417E60

Obfuscation Techniques

- Layout obfuscation is essentially useless.
- Renaming variables, removing whitespace and using #define routines for functions typically has very little impact on the underlying program.
- Sure you can do layout obfuscation on your code, and some of it MAY translate to obfuscated code, but the signal-to-noise ratio is much too low for to be useful.

MOV DWORD PTR DS: LECX+5C MOV DWORD PTR DS: LECX+5C MOV ECX.DWORD PTR SS: LEED MOV ECX.DWORD PTR SS: LEED MOV DWORD PTR DS: LECX+641 MOV EDX.DWORD PTR SS: LEED MOV EDX.DWORD PTR SS: LEED MOV EDX.DWORD PTR SS: LEED MOV ECX,DWORD PTR DS: L-8MSVCR90.PC411

JNZ

CAL

MOV ECX. DWORD PTR DS: LEDX MOV ECX. DWORD PTR SS: PUSH 00

Control-Flow Obfuscation

Turn that boring linear NOP sled into something worthy of Raging Waters.

Control-Flow Obfuscation

 With function pointers, method pointers, the volatile keyword and the goto keyword on our side, we can do some really fun stuff.

CL.B)

CAL

Control-Flow Obfuscation

- Opaque predicates are tautological IF statements.
- An opaque predicate cannot be optimized because the compiler cannot determine the outcome.
- You see this frequently in obfuscated JavaScript.

Control-Flow Obfuscation

int a=7,b=2,c=8,d=9;
if (a+b+c*d > 0)
{
 puts("yes");
 exit(0);
}
puts("no");

Control-Flow Obfuscation

int a=7,b=2,c=8,d=9;
if (a+b+c*d > 0)
{
 puts("yes");
 exit(0);
}
puts("no");

PUSH "yes" CALL \$PUTS PUSH Ø CALL \$EXIT

CALL 00401028 POP ECX POP ECX PUSH 195 PUSH 00417FF8 MOV ESI,00417E60

Control-Flow Obfuscation

int a,b,c,d; srand(time(0)); a=rand()+1;b=rand()+1; c=rand()+1; d=rand()+1;if (a+b+c*d > 0)puts("yes"); exit(0); puts("no");

> PUSH 00417FF8 MOV ESI,00417E60

Control-Flow Obfuscation

int a,b,c,d; srand(time(0)); a=rand()+1;b=rand()+1;c=rand()+1;d=rand()+1; if (a+b+c*d > 0)puts("yes"); exit(0); puts("no");

TEST EAX, EAX JLE SHORT :NO PUSH "yes" CALL \$PUTS PUSH 0 CALL \$EXIT NO: PUSH "no" CALL \$PUTS

PUSH 195 PUSH 00417FF8 MOV ESI,00417E60

Control-Flow Obfuscation

- Control-flow flattening involves, quite literally, flattening the graphical representation of your program.
- Typically you have a top-down flow with program graphs. With flattening, you cause a central piece of code to control the flow of the program.
- Control-flow obfuscation is employed by bin/crackmes/leetkey.exe

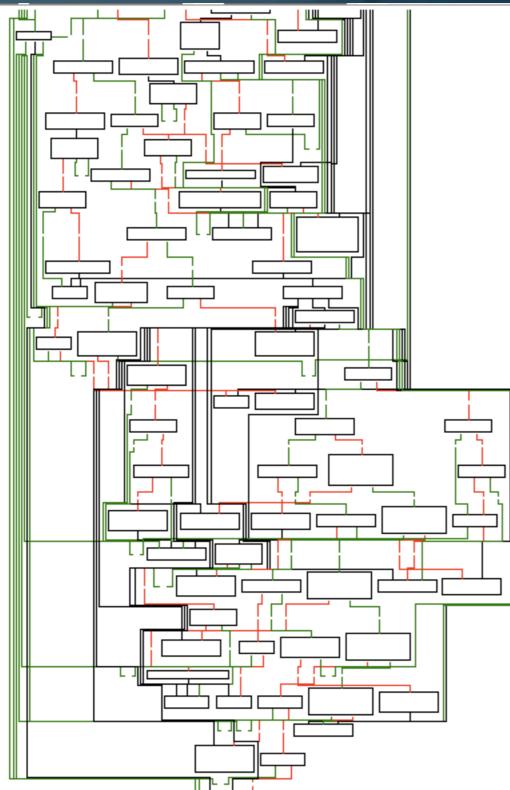
POP ECX PUSH 195 PUSH 00417FF8 MOV ESI 00417E60

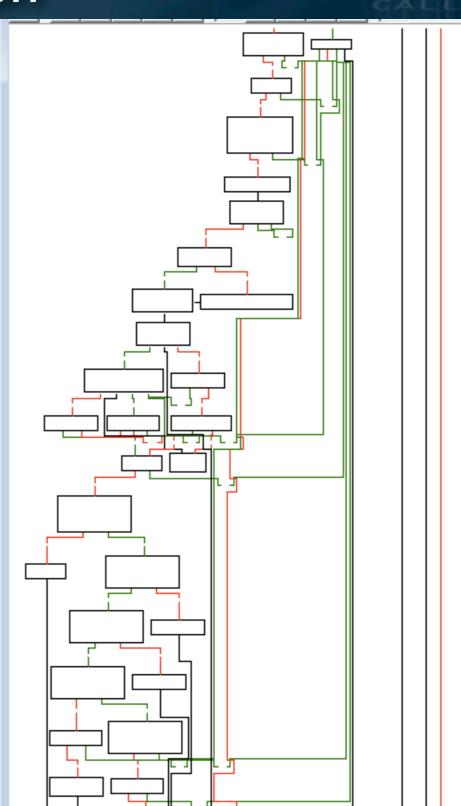
Control-Flow Obfuscation

Flattened:

Normal:

Control-Flow Obfuscation





Control-Flow Obfuscation

doThis(); doThat(); doMore();



int x=2;sw: switch(x) { case 0: doThat(); x = 1;goto sw; case 1: doMore(); break; case 2: doThis(); X = 0;goto sw;

> PUSH 00417FF8 MOV ESI,00417E60

Control-Flow Obfuscation

• This technique of obfuscation can be applied very creatively.

LEA EDX. DWOR

See src/cflow-flatlist.c and src/cflow-flattree.c

Control-Flow Obfuscation

- Most programs are reducible-- meaning they can easily be optimized.
- If a program is irreducible, then it cannot be optimized, thus translating spaghetti code into spaghetti assembly.
- A good example by Madou et. al. is making a loop irreducible.
- See src/cflow-irreducible.c

Control-Flow Obfuscation

- Raising bogus exceptions is a common way for malware to obfuscate and frustrate reverse engineering.
- This is easily accomplished by setting up a try block, intentionally triggering the exception, then resuming at the caught section.
- For Linux, you can do the same with signals.
- See src/cflow-exceptions.cpp

Control-Flow Obfuscation

try { volatile int trigger=20; doThis(); doThat(); /* trigger divide-by-zero exception */ trigger=trigger/(trigger-trigger); neverExecutes(); } catch (...) { doMore(); doTonsMore();

Data Obfuscation DWORD

ESTIDWORD

LEA EDX. DWORD JNZ SHC CALL 0

Ø

Data Obfuscation

- Data obfuscation takes a little more care than control-flow obfuscation.
- The data must be obfuscated before the compilation process, then de-obfuscated at run-time.
- If the data is not obfuscated before runtime, dead-code analysis is made trivial and your obfuscation is useless.

Data Obfuscation

- One of the more obvious techniques is to encrypt your strings.
- Even though strings don't technically lead to knowledge of the program, it can help aide in reverse-engineering more often than you think.

Data Obfuscation

- Recall the explanation of volatile:
 - volatile int x; x = 7; x <<= 2; x *= 2; x *= 2; x -= 12; x += (x*x)<<2; printf("%d\n", x);

• With enough annoyances, this can be used to frustrate analysis.

Data Obfuscation

• Data aggregation can be used to make dead-code analysis confusing.

char aggr[7] = "fboaor"; char foo[3], bar[3]; int i; for (i=0;i<3;++i) { foo[i]=aggr[i*2]; bar[i]=aggr[i*2+1]; } /* foo = "foo" / bar = "bar" *

Data Obfuscation

- Functions in the PLT/IAT are certainly considered data.
- To prevent dead-code analysis from discovering our library calls, we can easily "create" functions at run-time by using system calls such as LOadLibrary and GetProcAddress (Windows) and dlopen and dlsym (Linux).
- See src/data-loadlib.c, src/data-dlopen.c and src/mdl.cpp

Poor Man's Packer

How to simulate a packer in a humorous manner.

PUSH EAX PUSH EST CALL 00401028 POP ECX POP ECX PUSH 195 PUSH 00417FF8 MOV ESI,00417E60

SH

Poor Man's Packer

- Combines control-flow and data obfuscation to cause all sorts of headaches.
- Revolves around compiling, copying data and applying function pointers to obfuscated or encrypted data.
- See bin/crackmes/manifest.exe
 - If you have problems with this binary, ask a DC949 member what the group motto is.

POP ECX POP ECX PUSH 195 PUSH 00417FF8 MOV ESI,00417E60

- Compile
- Disassemble
- Copy bytes of function, make an array
- Apply encryption, aggregation, etc.
- Recompile
- Decipher at run-time
- Cast as function-pointer
- Execute
- See src/pmp-concept.c

- Problems
 - Functions are broken because they are no longer in the PLT/IAT.
 - Data offsets are completely messed up.
 - Functions in C++ objects cause segmentation faults (due to broken thiscall).
 - Compiler might change calling conventions.
 - void pointers are scary.

Poor Man's Packer

- If you pass a data structure containing data required by the function (function offsets, strings, etc.), you can circumvent the issue caused by relative jumps and offsets.
- This also applies to method pointers and C++ objects.

 This gives you the opportunity to dynamically add and remove necessary program data as you see fit.

- Be sure your calling conventions match after each step of compilation and bytecopying!
- cdecl is the calling convention used by vararg functions such as printf.
- fastcall and stdcall should be fine for all other functions.
- Mismatched calling conventions will cause headaches and segmentation faults.

- Why is this beneficial?
 - Ultimate control of all data
 - Code is still portable and executable
 - Adds a bizarre layer of obfuscation
 - When done enough, severely obfuscates source

- Why does this suck?
 - Makes binaries huge if you don't compress your functions due to enlarged data-sections
 - Takes a lot of work to accomplish
 - It can be extremely frustrating to craft the write code with the right keywords with full optimization

Additional Info

Some stuff to help you out with obfuscation

JNZ

SH

CALL 00401028 POP ECX POP ECX PUSH 195 PUSH 00417FF8 MOV ESI,00417E60

Tools

- Code transformers
 TXL (txl.ca)
 SUIF (suif.standford.edu)
- TXL and SUIF are used to transform source-code by a certain set of given rules (such as regular expressions).

Sources

- M. Madou, B. Anckaert, B. De Bus, K. De Bosschere, J. Cappaert, and B. Preneel, "On the Effectiveness of Source Code Transformations for Binary Obfuscation"
 B. M. Presed T. Chiuch, "A Binary Powriting
- B. M. Prasad, T. Chiueh, "A Binary Rewriting Defense against Stack based Buffer Overflows"

 C. I. Popov, S. Debray, G. Andrews, "Binary Obfuscation Using Signals" The End ESILDWORD LEA EDX. DWORD CL, BY JNZ SHO CALL ΡU SH004 ESI,00 0 MC